# Low-rank and sparse methods for high-dimensional approximation and model order reduction 

Lecture 3
Tensors - Sparse tensors

## Outline

(1) What are tensors?
(2) Sparse tensors
(3) Algorithms for sparse tensor approximation
(4) Sparse interpolation

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(1) What are tensors?
(2) Sparse tensors
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## Tensor product of vectors

For $I=\{1, \ldots, N\}$, an element $v$ of the vector space $\mathbb{R}^{I}$ is identified with the set of its coefficients $\left(v_{i}\right)_{i \in I}$ on a certain basis $\left\{e_{i}\right\}_{i \in I}$ of $\mathbb{R}^{\prime}$,

$$
v=\sum_{i \in I} v_{i} e_{i}
$$

Given $d$ index sets $I^{\nu}=\left\{1, \ldots, N_{\nu}\right\}, 1 \leq \nu \leq d$, we introduce the multi-index set

$$
I=I_{1} \times \ldots \times I_{d}
$$

An element $v$ of $\mathbb{R}^{/}$is called a tensor of order $d$ and is identified with a multidimensional array

$$
\left(v_{i}\right)_{i \in I}=\left(v_{i_{\mathbf{1}}}, \ldots, i_{d}\right)_{i_{\mathbf{1}} \in I_{\mathbf{1}}, \ldots, i_{d} \in I_{d}}
$$

which represents the coefficients of $v$ on a certain basis of $\mathbb{R}^{\prime}$.


## Tensor product of vectors

The entries of the multidimensional array are equivalently denoted

$$
v(i)=v\left(i_{1}, \ldots, i_{d}\right)
$$

Given $d$ vectors $v^{(\nu)} \in \mathbb{R}^{L_{\nu}}, 1 \leq \nu \leq d$, the tensor product of these vectors

$$
v:=v^{(1)} \otimes \ldots \otimes v^{(d)}
$$

is defined by

$$
v(i)=v^{(1)}\left(i_{1}\right) \ldots v^{(d)}\left(i_{d}\right)
$$

Such a tensor is called an elementary tensor. For $d=2$, using matrix notations, $v \otimes w$ is identified with the matrix $v w^{\top}$.


## Tensor product of vectors

The tensor space $\mathbb{R}^{I}=\mathbb{R}^{/_{\mathbf{1}} \times \ldots \times I_{d}}$, also denoted $\mathbb{R}^{/_{\mathbf{1}}} \otimes \ldots \otimes \mathbb{R}^{I_{d}}$, is defined by

$$
\mathbb{R}^{\prime}=\mathbb{R}^{/_{1}} \otimes \ldots \otimes \mathbb{R}^{I_{d}}=\operatorname{span}\left\{v^{(1)} \otimes \ldots \otimes v^{(d)}: v^{(\nu)} \in \mathbb{R}^{I^{\nu}}, 1 \leq \nu \leq d\right\}
$$

A canonical basis for $\mathbb{R}^{l}$ is given by

$$
e_{i}=e_{i_{1}} \otimes \ldots \otimes e_{i_{d}}, \quad i=\left(i_{1}, \ldots, i_{d}\right) \in I .
$$

where

$$
e_{i}(j)=e_{i_{1}}\left(j_{1}\right) \ldots e_{i_{d}}\left(j_{d}\right)=\delta_{i_{1}, j_{1}} \ldots \delta_{i_{d}, j_{d}}=\delta_{i, j}, \quad \forall j \in I
$$

A tensor $v \in \mathbb{R}^{\prime}$ can be written

$$
v=\sum_{i \in I} v(i) e_{i}=\sum_{i_{\mathbf{1}} \in I_{\mathbf{1}}} \ldots \sum_{i_{d} \in I_{d}} v\left(i_{1}, \ldots, i_{d}\right) e_{i_{\mathbf{1}}} \otimes \ldots \otimes e_{i_{d}}
$$

A natural norm on $\mathbb{R}^{\prime}$ is given by

$$
\|v\|=\sqrt{\sum_{i \in 1} v(i)^{2}}
$$

which makes $\mathbb{R}^{\prime}$ a Hilbert space.

## Tensor product of functions

Let $\mathcal{X}_{\nu} \subset \mathbb{R}, 1 \leq \nu \leq d$, be an interval and $V_{\nu}$ be a space of functions defined on $\mathcal{X}_{\nu}$. The tensor product of functions $v^{(\nu)} \in V_{\nu}$, denoted

$$
v=v^{(1)} \otimes \ldots \otimes v^{(d)}
$$

is a multivariate function defined on $\mathcal{X}=\mathcal{X}_{1} \times \ldots \times \mathcal{X}_{d}$ and such that

$$
v(x)=v\left(x_{1}, \ldots, x_{d}\right)=v^{(1)}\left(x_{1}\right) \ldots v^{(d)}\left(x_{d}\right)
$$

for $x=\left(x_{1}, \ldots, x_{d}\right) \in \mathcal{X}$. For example, for $i \in \mathbb{N}_{0}^{d}$, the monomial $x^{i}=x_{1}^{i_{1}} \ldots x_{d}^{i_{d}}$ is an elementary tensor.

## Tensor product of functions

The algebraic tensor product of spaces $V_{\nu}$ is defined as

$$
V_{1} \otimes \ldots \otimes V_{d}=\operatorname{span}\left\{v^{(1)} \otimes \ldots \otimes v^{(d)}: v^{(\nu)} \in V_{\nu}, 1 \leq \nu \leq d\right\}
$$

which is the space of multivariate functions $v$ which can be written as a finite linear combination of elementary (separated functions), i.e.

$$
v(x)=\sum_{k=1}^{n} v_{k}^{(1)}\left(x_{1}\right) \ldots v_{k}^{(d)}\left(x_{d}\right)
$$

Up to a formal definition of the tensor product $\otimes$, the above construction can be extended to arbitrary vector spaces $V_{\nu}$ (not only spaces of functions).

## Infinite dimensional tensor spaces

For infinite dimensional spaces $V_{\nu}$, a Hilbert (or Banach) tensor space equipped with a norm $\|\cdot\|$ is obtained by the completion (w.r.t. $\|\cdot\|$ ) of the algebraic tensor space

$$
\bar{V}^{\|\cdot\|}=\overline{V_{1} \otimes \ldots \otimes V_{d}}{ }^{\|\cdot\|}
$$

## Example 1 ( $L^{p}$ spaces)

Let $1 \leq p<\infty$. If $V_{\nu}=L_{\mu_{\nu}}^{p}\left(\mathcal{X}_{\nu}\right)$, then

$$
L_{\mu_{1}}^{p}\left(\mathcal{X}_{1}\right) \otimes \ldots \otimes L_{\mu_{d}}^{p}\left(\mathcal{X}_{d}\right) \subset L_{\mu}^{p}\left(\mathcal{X}_{1} \times \ldots \times \mathcal{X}_{d}\right)
$$

with $\mu=\mu_{1} \otimes \ldots \otimes \mu_{d}$, and

$$
\overline{L_{\mu_{1}}^{p}\left(\mathcal{X}_{1}\right) \otimes \ldots \otimes L_{\mu_{d}}^{p}\left(\mathcal{X}_{d}\right)}{ }^{\|\cdot\|}=L_{\mu}^{p}\left(\mathcal{X}_{1} \times \ldots \times \mathcal{X}_{d}\right)
$$

where $\|\cdot\|$ is the natural norm on $L_{\mu}^{p}\left(\mathcal{X}_{1} \times \ldots \times \mathcal{X}_{d}\right)$.

## Example 2 (Bochner spaces)

Let $\mathcal{X}$ be equipped with a finite measure $\mu$, and let $W$ be a Hilbert (or Banach) space. For $1 \leq p<\infty$, the Bochner space $L_{\mu}^{p}(\mathcal{X} ; W)$ is the set of Bochner-measurable functions $u: \mathcal{X} \rightarrow W$ with bounded norm $\|u\|_{p}=\left(\int_{\mathcal{X}}\|u(x)\|_{W}^{p} \mu(d x)\right)^{1 / p}$, and

$$
L_{\mu}^{p}(\mathcal{X} ; W)=\overline{W \otimes L_{\mu}^{p}(\mathcal{X})}{ }^{\|\cdot\|_{p}}
$$

## Tensor product basis

If $\left\{\psi_{i}^{(\nu)}\right\}_{i \in I_{\nu}}$ is a basis of $V_{\nu}$, then a basis of $V=V_{1} \otimes \ldots \otimes V_{d}$ is given by

$$
\left\{\psi_{i}=\psi_{i_{1}}^{(1)} \otimes \ldots \otimes \psi_{i_{d}}^{(d)}: i \in I=I_{1} \times \ldots \times I_{d}\right\}
$$

A tensor $v \in V$ admits a decomposition

$$
v=\sum_{i \in l} a(i) \psi_{i}=\sum_{i_{\mathbf{1}} \in I_{\mathbf{1}}} \ldots \sum_{i_{d} \in I_{d}} a\left(i_{1}, \ldots, i_{d}\right) \psi_{i_{\mathbf{1}}}^{(1)} \otimes \ldots \otimes \psi_{i_{d}}^{(d)}
$$

and $v$ can be identified with the set of its coefficients

$$
a \in \mathbb{R}^{I}
$$

## Orthogonal tensor product basis

If the $V_{\nu}$ are Hilbert spaces with inner products $(\cdot, \cdot)_{\nu}$ and associated norms $\|\cdot\|_{\nu}$, a canonical inner product on $V$ can be first defined for elementary tensors

$$
\left(v^{(1)} \otimes \ldots \otimes v^{(d)}, w^{(1)} \otimes \ldots \otimes w^{(d)}\right)=\left(v^{(1)}, w^{(1)}\right) \ldots\left(v^{(d)}, w^{(d)}\right)
$$

and then extended by linearity to the whole space $V$. The associated norm $\|\cdot\|$ is called the canonical norm.

If the $\left\{\psi_{i}^{(\nu)}\right\}_{i \in I_{\nu}}$ are orthonormal bases of spaces $V_{\nu}$, then $\left\{\psi_{i}\right\}_{i \in I}$ is an orthonormal basis of $\bar{V}^{\|\cdot\|}$. A tensor

$$
v=\sum_{i \in I} a_{i} \psi_{i}
$$

is such that

$$
\|v\|=\sqrt{\sum_{i \in 1} a_{i}^{2}}:=\|a\| .
$$

Therefore, the map $\psi$ which associates to a tensor $a \in \mathbb{R}^{\prime}$ the tensor $v=\Psi(a):=\sum_{i \in I} a_{i} \psi_{i}$ defines a linear isometry from $\mathbb{R}^{\prime}$ to $V$ for finite dimensional spaces, and between $\ell_{2}(I)$ and $\bar{V}^{\|\cdot\|}$ for infinite dimensional spaces.

## Curse of dimensionality

A tensor $a \in \mathbb{R}^{\prime}=\mathbb{R}^{I_{1} \times \ldots \times I_{d}}$ or a corresponding tensor $v=\sum_{i \in I} a_{i} \psi_{i}$, when $\# I_{\nu}=O(n)$ for each $\nu$, has a storage complexity

$$
\# I=\# I_{1} \ldots \# I_{d}=O\left(n^{d}\right)
$$

which grows exponentially with the dimension.

Manipulating tensors requires exploiting special properties of tensors such as sparsity or low rankness.

## Sparse tensors

## Outline

(2) What are tensors?
(2) Sparse tensors
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## Sparse tensors

A tensor $a \in \mathbb{R}^{I}$ is said to be $\Lambda$-sparse, with $\Lambda$ a subset of indices in $I$, if

$$
a_{i}=0 \quad \forall i \notin \Lambda .
$$

## A $\wedge$-sparse tensor (left) and the set $\wedge$ (right)




A tensor

$$
v=\sum_{i \in I} a_{i} \psi_{i} \in V
$$

is said to be $\Lambda$-sparse relatively to the tensor product basis $\left\{\psi_{i}=\psi_{i_{1}}^{(1)} \otimes \ldots \otimes \psi_{i_{d}}^{(d)}\right\}_{i \in I}$ of $V$ if the tensor $a=\left(a_{i}\right)_{i \in I}$ of its coefficients is $\Lambda$-sparse.

## Structured subsets of multi-indices

The set of multi-indices in $I=\mathbb{N}_{o}^{d}$ with $p$-norm bounded by $m$ is

$$
\Lambda=\left\{i \in I:\|i\|_{p} \leq m\right\}, \quad\|i\|_{p}= \begin{cases}\left(\sum_{k} i_{k}^{p}\right)^{1 / p} & \text { for } 0<p<\infty \\ \max _{k} i_{k} & \text { for } p=\infty\end{cases}
$$



## Structured subsets of multi-indices

## Example 3

Consider the tensor basis of monomials

$$
\psi_{i}(x)=x^{i}=x_{1}^{i_{1}} \ldots x_{d}^{i_{d}}
$$

- $\Lambda=\left\{i:\|i\|_{\infty} \leq m\right\}$ corresponds to the set of polynomials with partial degree less than $m$, with

$$
\# \Lambda=(m+1)^{d}
$$

- $\Lambda=\left\{i:\|i\|_{1} \leq m\right\}$ corresponds to the set of polynomials with total degree less than $m$, with

$$
\# \Lambda=\frac{(m+d)!}{m!d!}
$$

## Structured subsets of multi-indices

The set of multi-indices in $I=\mathbb{N}_{0}^{d}$ with weighted $p$-norm bounded by $m$ is

$$
\Lambda=\left\{i:\|i\|_{p, w} \leq m\right\}, \quad\|i\|_{p, w}= \begin{cases}\left(\sum_{k} w_{k} p i_{k}^{p}\right)^{1 / p} & \text { for } 0<p<\infty \\ \max _{k} w_{k} i_{k} & \text { for } p=\infty\end{cases}
$$

$$
m=6, p=\infty, w=(1,2) \quad m=6, p=1, w=(1,2) \quad m=6, p=1 / 2, w=(1,2)
$$





The choice of anisotropic sets is the key for high-dimensional approximation.
(3) Algorithms for sparse tensor approximation

- Greedy algorithms
- Convex relaxation methods
- Working set algorithms
(4) Sparse interpolation
(1) What are tensors ?

2 Sparse tensors
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## Greedy algorithms

Standard greedy algorithms can be applied for sparse tensor approximation by considering as the dictionary $\mathcal{D}$ the tensor product basis

$$
\mathcal{D}=\left\{\psi_{i}=\psi_{i_{1}}^{(1)} \otimes \ldots \otimes \psi_{i_{d}}^{(d)}: i \in I\right\} .
$$

Greedy algorithms construct a sequence of approximations $\left(u_{n}\right)_{n \geq 1}$ such that

$$
u_{n} \in \operatorname{span}\left\{\psi_{i}: i \in \Lambda_{n}\right\}
$$

where $\left(\Lambda_{n}\right)_{n \geq 1}$ is an increasing sequence of subsets such that

$$
\Lambda_{n}=\Lambda_{n-1} \cup\left\{i^{n}\right\}, \quad i^{n} \in I
$$


(3) Algorithms for sparse tensor approximation

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- Working set algorithms

4 Sparse interpolation

## Convex relaxation methods

Convex relaxation methods for sparse tensor approximation consists in solving problems of the form

$$
\min _{a \in \mathbb{R}^{\prime}} J(a)+\lambda \Omega(a)
$$

where $J$ is the cost functional and where $\Omega(a)$ is a convex function which induces sparsity in the resulting tensor $a \in \mathbb{R}^{\prime}$.

Note that in practice, for the problem to be tractable, the tensor is constrained to be $\Lambda$-sparse for a given $\Lambda$ with moderate cardinality.

## Convex relaxation methods

The standard approach uses the 1-norm

$$
\Omega(a)=\|a\|_{1}=\sum_{i \in I}\left|a_{i}\right|,
$$

which is the convex relaxation of the natural measure of sparsity $\|a\|_{0}=\#\left\{i \in I: a_{i} \neq 0\right\}$.
A weighted version of the 1 -norm can be defined as

$$
\Omega(a)=\|a\|_{1, \omega}=\sum_{i \in \Lambda} \omega_{i}\left|a_{i}\right|,
$$

where the weights can be deduced from a sequence of weights $\omega_{i_{\nu}}^{(\nu)}$ associated with univariate functions $\psi_{i_{\nu}}^{(\nu)}$ according to $\omega_{i}=\omega_{i_{1}}^{(1)} \ldots \omega_{i_{d}}^{(d)}$.

Also, by introducing a partition $\left\{\Lambda_{k}\right\}_{k=1}^{K}$ of $\Lambda$ (e.g. associated with a hierarchy of tensor spaces), we can consider

$$
\Omega(a)=\sum_{k=1}^{K} \omega_{k}\left\|a_{\wedge_{k}}\right\|_{2}
$$

for inducing group sparsity.
(1) What are tensors ?

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## Working set algorithms

Working set algorithms for sparse tensor approximation construct an increasing sequence of subsets $\left(\Lambda_{n}\right)_{n \geq 1}$ in $I$ and a sequence of approximations $u_{n} \in V_{\Lambda_{n}}$ computed through interpolation, regression or other projection methods.

The sequence of subsets is defined by

$$
\Lambda_{n}=\Lambda_{n-1} \cup A_{n}
$$

where $A_{n}$ is a subset of a candidate set $N_{n}$.



The definition of $N_{n}$ requires to define a strategy for the exploration of the set of multi-indices $I$.

## Hierarchy of tensor spaces

A $\Lambda$-sparse tensor $v$ is an element of the space

$$
V_{\Lambda}=\operatorname{span}\left\{\psi_{i}: i \in \Lambda\right\}
$$

A basis $\left\{\psi_{i}^{\nu}\right\}_{i \in I_{\nu}}$ in $V_{\nu}$ defines a hierarchy of spaces $H_{m}^{\nu}=\operatorname{span}\left\{\psi_{0}^{\nu}, \ldots, \psi_{m}^{\nu}\right\}$ such that

$$
\begin{equation*}
H_{0}^{\nu} \subset \ldots \subset H_{m}^{\nu} \subset \ldots \tag{1}
\end{equation*}
$$

For example, when considering canonical polynomial bases on $\mathcal{X}_{\nu}, H_{m}^{\nu}=\mathbb{P}_{m}\left(\mathcal{X}_{\nu}\right)$.
For $i \in I$, the product set

$$
R_{i}=\left\{0, \ldots i_{1}\right\} \times \ldots \times\left\{0, \ldots, i_{d}\right\}=\{j \in I: j \leq i\}
$$

defines a tensor space

$$
H_{i}:=V_{R_{i}}=H_{i_{1}}^{1} \otimes \ldots \otimes H_{i_{d}}^{d}
$$

which only depends on the spaces $H_{i_{\nu}}^{\nu}$ and not on the choice of bases $\left\{\psi_{0}^{\nu}, \ldots, \psi_{i_{\nu}}^{\nu}\right\}$ for $H_{i_{\nu}}^{\nu}$.
For two multi-indices $i, j \in I$,

$$
i \leq j \quad \text { implies } \quad R_{i} \subset R_{j} \quad \text { and } \quad H_{i} \subset H_{j}
$$

## Downward closed sets

A set $\Lambda \subset I$ is downward closed if for $j \in I$,

$$
(i \in \Lambda \text { and } j \leq i) \quad \Rightarrow \quad j \in \Lambda
$$



Not downward closed


## Downward closed sets

For $\Lambda$ downward closed, the space

$$
V_{\Lambda}=\operatorname{span}\left\{\psi_{i}: i \in \Lambda\right\}
$$

only depends on the hierarchy of spaces $H_{m}^{\nu}$ and not on the particular choice of bases $\left\{\psi_{0}^{\nu}, \ldots, \psi_{m}^{\nu}\right\}$ for these spaces.

## Example 4 (Polynomial spaces)

If $H_{m}^{k}=\mathbb{P}_{m}$ and $\Lambda$ is downward closed, then

$$
V_{\Lambda}=\operatorname{span}\left\{x^{i}=x_{1}^{i_{1}} \ldots x_{d}^{i_{d}}: i \in \Lambda\right\}
$$

whatever the choice of polynomial basis $\left\{\psi_{0}^{k}, \ldots, \psi_{m}^{k}\right\}$ of $\mathbb{P}_{m}$.

## Downward closed sets

An element $i$ of a set $\Lambda$ is said maximal if and only if there is no $j \in \Lambda$ such that $i \leq j$ and $i \neq j$.
A downward closed set $\Lambda$ is completely determined by the set of its maximal elements

$$
\Lambda=\bigcup_{\substack{i \in \Lambda \\ i \text { maximal }}} R_{i}
$$





The resulting space $V_{\wedge}$ is such that

$$
V_{\Lambda}=\sum_{\substack{i \in \Lambda \\ i \text { maximal }}} H_{i}=\sum_{\substack{i \in \Lambda \\ i \text { maximal }}} H_{i_{\mathbf{1}}}^{1} \otimes \ldots \otimes H_{i_{d}}^{d}
$$

## Margins

For a given set $\Lambda$, a natural neighborhood is the margin of $\Lambda$, defined by

$$
\mathcal{M}(\Lambda)=\left\{i \in I \backslash \Lambda: \exists j \in \Lambda \text { s.t. }\|i-j\|_{1}=1\right\}
$$

Also, we define the reduced margin of $\Lambda$ as

$$
\mathcal{M}_{r}(\Lambda)=\left\{i \in I \backslash \Lambda: i-e_{k} \in \Lambda \text { for all } k \text { s.t. } i_{k}>1\right\}
$$

## A set $\Lambda$ and its margin $\mathcal{M}(\Lambda)$



## A set $\Lambda$ and its reduced margin $\mathcal{M}_{r}(\Lambda)$



For a downward closed set $\Lambda$, an interesting property of the reduced margin $\mathcal{M}_{r}(\Lambda)$ is that for any subset $A \subset \mathcal{M}_{r}(\Lambda), \Lambda \cup A$ is downward closed.

## Sparse interpolation

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## Sparse interpolation

## Sparse interpolation

Let $\left\{t_{i}^{\nu}\right\}_{i \geq 0}$ be a sequence of points such that $\Gamma_{m}^{\nu}=\left\{t_{0}^{\nu}, \ldots, t_{m}^{\nu}\right\}$ is unisolvent for $H_{m}^{\nu}$, that means for all $y \in \mathbb{R}^{m+1}$, there exists a unique function $v \in H_{m}^{\nu}$ such that $v\left(t_{i}^{\nu}\right)=y_{i}$ for all $0 \leq i \leq m$. Let $\left\{\ell_{i}^{\nu, m}(t)\right\}_{i=0}^{m}$ be the interpolation functions in $H_{m}^{\nu}$.

We define the interpolation operator $\mathcal{I}_{m}^{\nu}$ which associates to a function $v$ the unique function $\mathcal{I}_{m}^{\nu}(v) \in H_{m}^{\nu}$ such that $\mathcal{I}_{m}^{\nu}(v)\left(t_{i}^{\nu}\right)=v\left(t_{i}^{\nu}\right)$ for $0 \leq i \leq m$,

$$
\mathcal{I}_{m}^{\nu}(v)(t)=\sum_{i=0}^{m} v\left(t_{i}^{\nu}\right) \ell_{i}^{\nu, m}(t)
$$

The interpolation operator can be written as the sum of difference operators

$$
\mathcal{I}_{m}^{\nu}=\sum_{i=0}^{m} \Delta_{i}^{\nu}
$$

with

$$
\Delta_{i}^{\nu}=\mathcal{I}_{i}^{\nu}-\mathcal{I}_{i-1}^{\nu}, \quad \mathcal{I}_{0}^{\nu}=0
$$

## Sparse interpolation

## Sparse interpolation

To a multi-index $i$ corresponds a point

$$
x_{i}=\left(t_{i_{1}}^{1}, \ldots, t_{i_{d}}^{d}\right)
$$

For the product set of indices $R_{i}=\{j \in I: j \leq i\}$, we define the tensor product grid

$$
\Gamma_{i}=\left\{x_{i}: i \in R_{i}\right\},
$$

which is unisolvent for the tensor space $V_{R_{i}}=H_{i}=H_{i_{1}}^{1} \otimes \ldots \otimes H_{i_{d}}^{d}$, and the associated interpolation operator $\mathcal{I}_{i}=\mathcal{I}_{i_{1}}^{1} \otimes \ldots \otimes \mathcal{I}_{i_{d}}^{d}$ such that

$$
\mathcal{I}_{i}(v)(x)=\sum_{j_{1}=0}^{i_{1}} \ldots \sum_{j_{d}=0}^{i_{d}} v\left(t_{j_{1}}^{1}, \ldots, t_{j_{d}}^{d}\right) \ell_{j_{1}}^{1, i_{1}}\left(x_{1}\right) \ldots \ell_{j_{d}}^{d, i_{d}}\left(x_{d}\right):=\sum_{j \in \Gamma_{i}} v\left(x_{j}\right) \ell_{j}^{i}(x)
$$

$\mathcal{I}_{i}$ can be expressed as

$$
\mathcal{I}_{i}=\sum_{j \leq i} \Delta_{j}
$$

where

$$
\Delta_{j}=\Delta_{j_{1}}^{1} \otimes \ldots \otimes \Delta_{j_{d}}^{d} .
$$

## Sparse interpolation

## Sparse interpolation

The above construction can be generalized to an arbitrary downward closed set $\Lambda$, with the definition of an interpolation operator

$$
\mathcal{I}_{\Lambda}=\sum_{i \in \Lambda} \Delta_{i}
$$

onto the space $V_{\Lambda}$, associated with the sparse grid

$$
\Gamma_{\wedge}=\left\{x_{i}: i \in \Lambda\right\} .
$$

The grid $\Gamma_{\Lambda}$ is unisolvent for the space $V_{\Lambda}$.

